

# ALMOST GROUP ENVY-FREE ALLOCATION OF INDIVISIBLE GOODS AND CHORES

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## Research Question

How to define fairness criteria for *groups of agents* when considering the allocation of indivisible *goods and chores*?

The main idea is to extend the *group fairness criteria* introduced by Conitzer et al. (2019) when there are only goods to the framework of fair division of *goods and chores* introduced by Aziz et al. (2019).

## Setting

We are interested in *Fair Division* problems where:

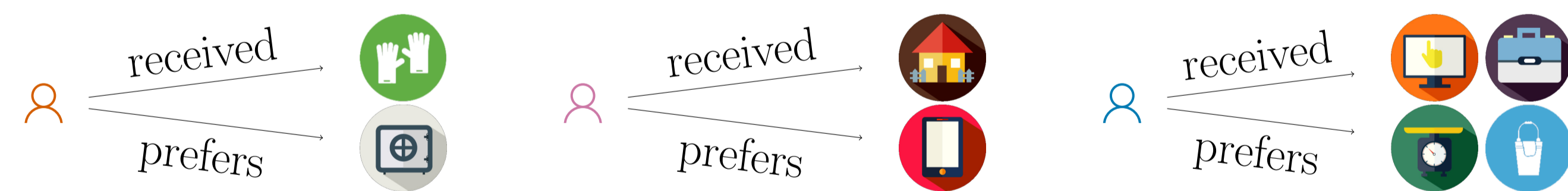
- Items are *indivisible*
- Items can be considered as *goods*  and/or *chores* 
- Agents have *additive* preferences

## Group envy-freeness

An allocation is GEF if there are *no* two groups  $S$  and  $T$  with  $|S| = |T|$  such that:



meaning that there is a *reallocation* of the items received by agents in  $T$  to agents  $S$  that *Pareto-dominates* the current allocation for agents in  $S$ :



## Theorems

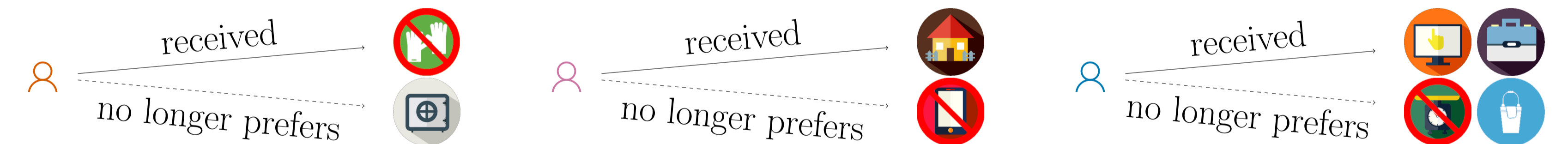
When agents have *additive* and *identical* preferences, a GEF1 allocation *always exists* and can be computed in *polynomial time*.

When agents have *additive* and *ternary symmetric*<sup>1</sup> preferences, a GEF1 allocation *always exists* and can be computed in *polynomial time*.

<sup>1</sup> For every agent, the preferences of the single items are always in  $\{\alpha, 0, -\alpha\}$ .

## Group envy-freeness up to one item

We can decide *not to consider an item* for each agent in  $S$  when checking whether she prefers the reallocation. If that eliminates the envy, the allocation is GEF1:



An allocation  $\pi$  is GEF1 if for every two groups  $S$  and  $T, |S| = |T| \neq 0$ ,

1. for every reallocation  $\pi'$  of the items in  $T$  to the agents in  $S$ ,
2. for every agent  $i \in S$ , there exists an item  $o_i$  in what she received either in the current allocation or in the reallocation such that,
3. by removing all  $o_i$ , the reallocation  $\pi'$  does not Pareto-dominate  $\pi$ .

## Theorem

Checking whether an allocation is GEF1 is *coNP-complete* for additive preferences and where there are only goods, only chores or both goods and chores.

Aziz, H.; Caragiannis, I.; Igarashi, A.; and Walsh, T. 2019. Fair Allocation of Indivisible Goods and Chores. In *Proc. of 28th IJCAI*.

Conitzer, V.; Freeman, R.; Shah, N.; and Vaughan, J. W. 2019. Group Fairness for the Allocation of Indivisible Goods. In *Proc. of 33rd AAAI Conference*.